Based on the Circular Logarithm Theory "Verification of Falk-Wills Theorem" and Its Application

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Abstract: Fermat's theorem has been dialectically verified by more than three hundred years of history and multi-personal conjecture. Unfortunately, there is only a difference between the inequality and the equation, and no integer expansion of compatibility is found. The conclusion is unfair. Defining group algebraic closed chains is an ordered combination set and equilibrium of finite finite-dimensional matrices of infinite elements, proving its unity, reciprocity, isomorphism, parallelism, limit theorem, and considering scalability, security, and decentralization. The superiority of chemistry, the logarithm of the dimensionless quantum function is established, and the arithmetic operation between [0~1] is realized, which is called the logarithm theory. It is convincingly proved that any inequality is an integer equality expansion, and it is proposed to prosecute Bug damage and solve the impossible triangle idea. The circular logarithm theory is proved by the "angel particle" of the physical experiment. In the actual project, there are innovative inventions for asymmetric energy such as earth electromagnetic field generators, gravity field engines, and turbojet engines.

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1. Introduction

Fermat Dading reason French mathematician Fermat asserts that when the integer P>2, the equation for A, B, C, $A^P + B^P = C^P$ has no positive integer solution. After being proposed, he experienced many people's conjectures and dialectic. After more than three hundred years of history, he was certified by British mathematician Andrew Wiles in 1995. Unfortunately, this proof only differs between the inequality and the equation, and no compatibility is found to achieve zero-error integer expansion. The conclusion is unfair, collectively known as the Fermat-Wills inequality theorem.

The controversial focus of the Fermat-Wills inequality theorem: the associated equations differ from the inequalities. Is there compatibility? How to use the compatibility to convert into a self-consistent integer equality expansion, and overcome the current hot blockchain chain defects, to achieve organic integration and complement each other.

This paper proposes the concept of group algebra closed chain. Define the state of the equilibrium movement of the group algebraic closed chain as a point set, also known as the "point state". Is an algebraic combination (±p) of an infinite element (Z) arbitrarily complex (±S±N) dimension (with calculus order ±N), and becomes an algebraic cluster of power functions (Z±S±N±p), both Integer (±S±N±p)=(±1) Infinitely ordered integer or simple paste expansion, proving its inversion, unity, reciprocity, isomorphism, parallelism, and limitless structural features; It also has the advantages of scalability, safety, and centering. Cause any inequality to be converted to a complete equality. False Falma-Wills Theorem.

Thus, there is a logarithm with a dimensionless quantity elliptic function as the base, called the circular logarithm algorithm (also known as the relativistic supersymmetric element matrix), which realizes the perfect combination of the circular log-block chain. The arithmetic operation of the probability quantum "expansion" and "no specific particle content" in [0~1].

2. Gains and losses of Wiles's theorem

Ribet proved in 1986 that the Frey curve does not have a mode. Encouraged by Ribet's work, Wiles spent six years trying to prove that each (or at least most of) elliptic curves have a modular pattern. In the end, he proved that each elliptic curve (an elliptic curve is semi-stable, and the prime p can be proved as the prime factor of the discriminant of E) has a modal curve; since the Frey curve is semi-stable, this is enough to derive the fee. Ma's theorem, there are:

The Wiles theorem tells us that there is a model mode, that is, the prime p-deficiency has a modular mode.

Ribet's theorem tells us that we have no modularity. Ribet proved in 1986 that the Frey curve does not have a mode. Encouraged by Ribet's work, Wiles spent six years trying to prove that each (or at least most of) elliptic curves have a modular pattern. In the end, he proved that each elliptic curve (an elliptic curve is semi-stable, and the prime p can be proved as the prime factor of the discriminant of E) has a modal curve; since the Frey curve is semi-stable, this is enough to derive the fee. Ma's theorem, there are:

The Wiles theorem tells us that there is a model mode, that is, the prime p-deficiency has a modular mode.

Ribet's theorem tells us that we have no modularity. At this point,

Wiles said: The above incompatibility leads to the equation without a non-zero integer solution, the equation and the inequality are not unified, and the Fermat's theorem is established. However, this proof only sees the difference between the inequality and the equation, does not find compatibility, and the conclusion is unfair.

In the exploration of the inequality and equality relationship of Fermat-Wills theorem, we find that the multiplication of uncertainty can be converted into the reciprocal average, and the unit of the positive mean has the inversion and compatibility. The inequality of uncertainty is convincingly transformed into an integer equality of zero error expansion.

(1), The result obtained by the Wiles theorem: the difference between the equation and the inequality is obtained, and the compatibility is not found. The conclusion is that it cannot be unified. The essence is whether the entanglement analysis and the discrete statistics can be unified.

have:

$$A^{K(Z\pm S\pm N\pm P)}+B^{K(Z\pm S\pm N\pm P)}\neq C^{K(Z\pm S\pm N\pm P)}; \qquad (1.1)$$

(2) The result obtained by the circular logarithm theorem: the difference between the equation and the inequality is obtained, and the compatibility is found. The conclusion is that it can be unified. The essence is that the entanglement analysis and the discrete statistics can be integrated into one.

have:

$$\{A\}^{K(Z\pm S\pm N\pm P)}+\{B\}^{K(Z\pm S\pm N\pm P)}$$
$$=(1-\eta^2)\{C\}^{K(Z\pm S\pm N\pm P)}; \qquad (1.2)$$
$$0 \leqslant [(1-\eta^2)\sim(\eta)]^{K(Z\pm S\pm N\pm P)} \leqslant 1; \qquad (1.3)$$

Where: {A}, {B}, {C} are all integers or prime numbers. Fermat's theorem and the Wiles's theorem inequality, which maintains the integer equality of C by the logarithm of the circle $(1-\eta^2)$. $(1-\eta^2)$ is a reciprocal equation of equality that is completed by circular logarithm reconstruction. It is called "topological quantum" in the blockchain.

Said above:

(1) Discrete mathematics: refers to "there is no interaction between elements in the element group". The change of one element in the group does not affect the change effect of the whole value, and satisfies the axiomatization of the set theory "self divided by itself equal to 1. Assume that you can satisfy the complete equality equation. The blockchain is called the "branch state".

(2) Entangled mathematics: refers to the "interaction between elements in the element group. When any one element changes, the other elements in the group change accordingly, affecting the overall effect of the group", and get "self divided by itself. The elliptic function topology and probability structure, which must be equal to 1", are the basis for establishing inequalities. In the blockchain, it is called "superimposed state".

3. "Reciprocal mean and positive mean" constitute an interactive inversion law

[Lemma 1] Multi-element multiplication is a reciprocal "reciprocal mean and positive mean"

Definition: Mean function value: the number of non-repetitive combinations of infinite elements in the closed chain of the infinite elements of the group algebra, except for the number of its corresponding combination forms (called coefficients) such as: P combination coefficient $C_{(S\pm P)}$ regularization conditions, $C_{(S+P)}=C_{(S-P)}$;

$$\{X_0\}^{K(Z\pm S\pm N\pm P)}=\{D_0\}^{K(Z\pm S\pm N\pm P)}$$
$$=\{\Sigma(1/C_{(S\pm P)})[\Pi_P X_i^{K}+\cdots]\}^{K(Z\pm S\pm N\pm P)}$$

Yes: $\{X\}^{K(Z\pm S\pm N-P)}$ in the unknown function is called

(P=-P) reciprocal function value;

$\{X_0\}^{K(Z\pm S\pm N-P)}$ in the unknown mean function is called (P=-p) the average of the reciprocal function;

$\{D_0\}^{K(Z\pm S\pm N+P)}$ in the unknown mean function is called (P=+p) the average of the reciprocal function;

The known function;

$\{X\pm D\}^{K(Z\pm S\pm N\pm P)}$ is called (P=±p) combination equation in the combination function;

$\{X_0\pm D_0\}^{K(Z\pm S\pm N\pm P)}$ in the combined averaging function (P=±p) combined averaging equation;

(Note: Sometimes the length of the province (±S±N)

or (±N) is not written, which means the general formula, the same below)

Assume

$$\{X\}^{K(Z\pm S-P)}$$
$$=\Pi(x_a^{-1}\cdot x_b^{-1}\cdot\cdots\cdot x_p^{-1}\cdot\cdots\cdot x_q^{-1})^{K(Z\pm S-P)};$$
$$\{D\}^{K(Z\pm S+P)}$$
$$=\Pi(D_a^{+1}\cdot D_b^{+1}\cdot\cdots\cdot D_p^{+1}\cdot\cdots\cdot D_q^{+1})^{K(Z\pm S+P)};$$
$$\{X_0\}^{K(Z\pm S-P)}=[(1/C_{(S-P)})^{-1}\Sigma(\Pi_P x_a^{-1}$$
$$+\Pi_P x_b^{-1}+\cdots+\Pi_P x_p^{-1}+\cdots+\Pi_P x_q^{-1})_S]^{K(Z\pm S-P)};$$
$$\{D_0\}^{K(Z\pm S+P)}=[(1/C_{(S+P)})^{+1}\Sigma(\Pi_P D_a^{+1}$$
$$+\Pi_P D_b^{+1}+\cdots+\Pi_P D_p^{+1}+\cdots+\Pi_P D_q^{+1})_S]^{K(Z\pm S+P)};$$
$$\{X\}^{K(Z\pm S-P)}+\{D\}^{+(Z\pm S+P)}=(1/2)\{X\pm D\}^{0(Z\pm S+P)};$$
$$(1-\eta^2)^{K(Z\pm S\pm P)}=\{X_0\}^{K(Z\pm S-P)}\cdot\{D_0\}^{K(Z\pm S+P)}$$

$= \{X_0\}^{K(Z\pm S\pm P)}/\{D_0\}^{K(Z\pm S\pm P)}$;

Prove: Interactive reversal of each combination (S=±P) level.certificate:

Take the iterative method of p=±1, and divide it by

$[(1/C_{p+1})(x_a+x_b+\cdots+x_p+\cdots+x_q)]^{K(Z\pm S+1)}$

$\{X\}^{K(Z\pm S\pm 1)}=[\prod(x_a\bullet x_b\bullet\cdots\bullet x_p\bullet\cdots\bullet x_q)^{K(Z\pm S\pm 1)}$

$=[(C_{p+1})\prod(x_a\bullet x_b\bullet\cdots\bullet x_p\bullet\cdots\bullet x_q)^{K(Z\pm S\pm 1)}$

$/(1/C_{p+1})(x_a+x_b+\cdots+x_p+\cdots+x_q)^{K(Z\pm S+1)}]$

$\bullet(1/C_{p+1})(x_a+x_b+\cdots+x_p+\cdots+x_q)^{K(Z\pm S+1)})$

$=[(1/C_{(P-1)})^{-1}\Sigma(x_a^{-1}+x_b^{-1}+\cdots+x_p^{-1}+\cdots+x_q^{-1})]^{K(Z\pm S-1)}$

$\bullet[(1/C_{(S+P)})^{+1}\Sigma(D_a^{+1}+D_b^{+1}+\cdots+D_p^{+1}+\cdots+D_q^{+1})]^{K(Z\pm S+1)}$

$=\{X_0\}^{K(Z\pm S-1)}\bullet\{D_0\}^{K(Z\pm S+1)}$ (2.1)

In the middle:

$\{X_0\}^{K(Z\pm S-1)}=\{D_0\}^{K(Z\pm S+1)}$

$=[(1/C_{(S+P)})^{+1}\Sigma(D_a^{+1}+D_b^{+1}+\cdots+D_p^{+1}+\cdots+D_q^{+1})]^{K(Z\pm S+1)}$

on the contrary:

$\{X\}^{K(Z\pm S\pm 1)}=[(C_{p+0})\prod_p(x_a\bullet x_b\bullet\cdots\bullet x_p\bullet\cdots\bullet x_q)]^{K(Z\pm S\pm 0)}$

$/[(1/C_{p-1})^{-1}(x_a^{-1}+x_b^{-1}+\cdots+x_p^{-1}+\cdots+x_q^{-1})]^{K(Z\pm S-1)}$

$\bullet[(1/C_{p-1})^{-1}(D_a^{-1}+D_b^{-1}+\cdots+D_p^{-1}+\cdots+D_q^{-1})]^{K(Z\pm S-1)}$

$=\{X_0\}^{K(Z\pm S-1)}\{D_0\}^{K(Z\pm S+1)}$ (2.2)

The same reason: can be analogized by order (P = 0, 1, 2, 3, 4,... natural number).

have:

$\{X\}^{K(Z\pm S\pm p)}$

$=[(1/C_{S\pm p})\prod(x_a\bullet x_b\bullet\cdots\bullet x_p\bullet\cdots\bullet x_q)]^{K(Z\pm S\pm p)}$

$/(1/C_{S+p})\Sigma(\prod x_a+\prod x_b+\cdots+\prod x_p+\cdots+\prod x_q)^{K(Z\pm S+p)}$

$\bullet(1/C_{S-P})\Sigma(\prod D_a+\prod D_b+\cdots+\prod D_p+\cdots+\prod D_q)^{K(Z\pm S-p)}$

$=\{X_0\}^{K(Z\pm S-p)}\bullet\{D_0\}^{K(Z\pm S+p)}$ (2.3)

equations (2.1)~(2.3) prove that any multi-element multiplication is essentially the combination of the "positive mean" and "reciprocal mean" of reciprocity.

where: any finite finite power power infinite is a set of Z=K (Z±S±P), (Z) represents the completeness of the infinite element of the algebraic closed chain, and (±S) represents any finite complex dimension in the closed set of groups., (±P) Algebraic clusters of $\{x\}^{K(Z\pm S\pm N\pm P)}$ without repeated combinations of all elements.

[Lemma 2] The logarithm of the circle reflects the rule of change between the "reciprocal mean and the positive mean"

Proof: Algebraic closed chain algebraic clusters have isomorphic reciprocal inversion

Further derivation according to formula (2.3):

heve:

$\{X\}^{K(Z\pm S-P)}$

$=\{X_0\}^{K(Z\pm S-P)}/\{D_0\}^{K(Z\pm S+P)}\bullet\{D_0\}^{K(Z\pm S+P)}$

$=(1-\eta^2)^{K(Z\pm S\pm P)}\{D_0\}^{K(Z\pm S+P)}$

$=\{0\sim1\}\{D_0\}^{K(Z\pm S+P)}$; (3.1)

among them:

$(1-\eta^2)^{K(Z\pm S\pm P)}$

$=\{X_0\}^{K(Z\pm S-P)}\bullet\{D_0\}^{K(Z\pm S+P)}$

$=[\{X_0\}/\{D_0\}]^{K(Z\pm S\pm P)}$; (3.2)

$0\leqslant(1-\eta^2)^{K(Z\pm S\pm P)}$

$=(1-\eta^2)^{K(Z\pm S+P)}\bullet(1-\eta^2)^{K(Z\pm S-P)}\leqslant\{1\}^{K(Z\pm S\pm P)}$; (3.3)

In the middl:

$\{X\}^{K(Z\pm S-P)}$

$=(1+\eta)\{X_0\}^{K(Z\pm S-P)}=(1-\eta^2)\{X_0\}^{K(Z\pm S-P)}$; (3.4)

$\{D\}^{K(Z\pm S+P)}$

$=(1-\eta)\{D_0\}^{K(Z\pm S+P)}=(1-\eta^2)\{D_0\}^{K(Z\pm S+P)}$; (3.5)

The merger is written as:

$W=(1-\eta^2)^Z W_0$; (3.6)

$0\leqslant[(1-\eta^2)\sim(\eta)]^Z\leqslant1$; (3.7)

Where: W, W_0 represent arbitrary unknown, known group set, algebraic closed chain, geometric space, numerical value, probability, topology, event. $(1-\eta^2)^Z$ represents the reciprocal change rule of each algebraic cluster of group elements, called the logarithm of the circle.

Any dimensional inequality is transformed into a balanced integer equality, and the unitary topological expansion is obtained, which gives a self-contained, unified all-element description. Produce the following effects:

(1) 、Replace the group theory "if and only" with completeness with the "=" symbol.

(2), "Arithmetic four arithmetic symbols" integrity calculation replaces "logical arithmetic symbols"

In particular, the natural rule of reciprocal "reciprocal mean" and "positive mean" has not been discovered before. Its appearance avoids the phenomenon of mathematical lameness, making mathematics more complete, complete and concise.

4. Algebraic closed chain has the isomorphism of isomorphic topology

It is now continued to prove that the group algebraic closed chain has isomorphic reciprocal dynamics under dynamic equilibrium and imbalance conditions, and its circular logarithm power function plus (/t) becomes a dynamic expression.

Let: inequalities and equations or unbalanced and balanced group algebra closed-chain dynamic equation $\{X \pm D\}^{k(Z \pm S \pm P)/t}$;

heve:

$\{X\}^{K(Z\pm S-p)/t} \neq \{D\}^{K(Z\pm S+p)/t}$;

or: $\{^{KS}\sqrt{D}\}^{K(Z\pm S+p)/t} \neq \{D_0\}^{K(Z\pm S+p)/t}$

$B=[(1/C_{S+1})(D_a+D_b+\cdots+D_P+\cdots+D_q)]^{k(Z\pm S\pm 1)/t}$
$=\{D_0\}^{k(Z\pm S\pm 1)/t}$,

$P=[(1/C_{S+P})^K(\prod_P D_a^K+\prod_P D_b^K+\cdots+\prod_P D_P^K+\cdots$
$+\prod_P D_q^K)]^{k(Z\pm S\pm P)/t}=\{D_0\}^{k(Z\pm S\pm P)/t}$,

Proof: In the polynomial regularization, the (second) coefficient B and (the last second term) Q are divided into a combination form, expressing the average of the function.

heve:

$\{X\}=\{^{KS}\sqrt{D}\}^{K(Z\pm S-p)/t}=\{D\}$,

$\{X_0\}^{K(Z\pm S-p)/t}=\{D_0\}^{K(Z\pm S+p)/t}$;

$(1-\eta^2)^{(Z/t)}=\{X\}^{K(Z\pm S+0)/t}/\{D\}^{K(Z\pm S+0)/t}$,

$=\{^{KS}\sqrt{D}\}^{k(Z\pm S+1)/t}/\{D_0\}^{k(Z\pm S+1)/t},\cdots$

$=\{^{KS}\sqrt{D}\}^{k(Z\pm S+p)/t}/\{D_0\}^{k(Z\pm S+p)/t},\cdots$

$=\{^{KS}\sqrt{D}\}^{k(Z\pm S+q)/t}/\{D_0\}^{k(Z\pm S+q)/t}$,

$$\{X\pm D\}^{(Z/t)}=Ax^{K(Z\pm S-0)/t}+Bx^{K(Z\pm S-1)/t}+\cdots+Px^{K(Z\pm S-p)/t}+\cdots+Qx^{K(Z\pm S-q)/t}+D^{K(Z\pm S+0)/t}$$
$$=C_{(S-0)}x^{K(Z\pm S-0)/t}\bullet D_0^{k(Z\pm S+0)/t}$$
$$+C_{(S-1)}x^{K(Z\pm S-1)/t}\bullet D_0^{k(Z\pm S+1)/t}+\cdots$$
$$+C_{(S-p)}x^{K(Z\pm S-p)/t}\bullet D_0^{k(Z\pm S+p)/t}+\cdots$$
$$+C_{(S-p)}x^{K(Z\pm S-q)/t}\bullet D_0^{k(Z\pm S+p)/t}\pm C_{(S+0)}D^{K(Z\pm S+0)/t}$$
$$=x_0^{K(Z\pm S-0)/t}D_0^{K(Z\pm S+0)/t}+x_0^{K(Z\pm S-1)/t}D_0^{K(Z\pm S+1)/t}+\cdots$$
$$+x_0^{K(Z\pm S-p)/t}D_0^{K(Z\pm S+p)/t}+\cdots$$
$$+x_0^{K(Z\pm S-q)/t}D_0^{K(Z\pm S+q)/t}\pm\{^{KS}\sqrt{D}\}^{+(Z\pm S+p)/t}$$

The two sides of the equal sign are divided by $\{X_0 \pm D_0\}^{K(Z\pm S)/t}$ to obtain the series expansion of $(1-\eta^2)^{K(Z\pm S)/t}$

get:

$\{X\}^{K(Z\pm S)/t}\pm\{D\}^{K(Z\pm S)/t}$
$=(1/2)\{X\pm D\}^{K(Z\pm S)/t}$
$=(1/2)[\{^{KS}\sqrt{D}\}\pm\{D_0\}]^{K(Z\pm S)/t}$
$=(1-\eta^2)^{K(Z\pm S)/t}\{0,1/2,1\}^{K(Z\pm S)/t}\{D_0\}^{K(Z\pm S)/t}$; (4.1)

$(1-\eta^2)^{(Z/t)}=(1-\eta^2)^{K(Z-0)/t}+(1-\eta^2)^{K(Z-1)/t}+\cdots$
$+(1-\eta^2)^{K(Z-p)/t}+\cdots+(1-\eta^2)^{K(Z-q)/t}$ (4.2)

$0\leqslant(1-\eta^2)^{(Z/t)}\sim(\eta)^{(Z/t)}\leqslant 1$; (4.3)

Where: $(1-\eta^2)^{(Z/t)}=\{0 \text{ or } 1\}^{(Z/t)}$ is a discrete statistic, a circular logarithmic limit (or center point/boundary condition); $0\leqslant(1-\eta^2)^{(Z/t)}\sim(\eta)^{(Z/t)}\leqslant 1$ is an entanglement analysis, that is, topology and probability conditions.

The formulas (4.1)~(4.3) prove that the total element dimension is invariant, even if it is an asymmetrical combination, the relative balance of each level is formed by the logarithm of the circle, the inequality is satisfied to become the equation, and the crisis of inequality and equality is solved., indicating that the inequality is more basic than the equation. In physics, it is called "chirality." Get mathematical proof here.

5. Three norm invariance theorems for circular logarithms

The reversal of the interaction between the "positive mean" and the "reciprocal mean". The integer zero error expansion of the unit cell of the group algebraic closed chain is satisfied to ensure the unity of its power function. However, in the unitary topology, there are three norm invariants and limits unique to the circular logarithm, which is an important theorem for the conversion of inequalities into inequalities.

[Theorem 1], the first norm invariance theorem (unit log logarithm):

Unit circle logarithm: "The sum of the set items of its own elements $\sum\{x_h\}^{(Z/t)}$ divided by the total set of its own elements $(1-\eta_H^2)^{(Z/t)}$ the sum must be equal to $\{1\}^{(Z/t)}$ ". That is, the multi-element "normalization" is the unit body group algebraic closed chain.

heve:

$(1-\eta_H^2)^{K(Z\pm S\pm P)}=\sum\{x_h\}^{(Z/t)}/\{x_H\}^{(Z/t)}$
$=\{[\Sigma(\prod x_{h1}+\prod x_{h2}+\cdots+\prod x_{hp}+\cdots+\prod x_{hq})]$
$/\{x_H\}\}^{K(Z\pm S\pm P)}$
$=\{(1-\eta_{h1}^2)+(1-\eta_{h2}^2)+\cdots+(1-\eta_p^2)+\cdots$
$+(1-\eta_q^2)\}^{K(Z\pm S\pm P)}$
$=\{1\}^{K(Z\pm S\pm P)}$; (5.1)

or:

$(\eta_H)^{K(Z\pm S\pm P)}=\{\eta_{h1}+\eta_{h2}+\cdots+\eta_p+\cdots+\eta_q\}^{K(Z\pm S\pm P)}$
$=\{1\}^{K(Z\pm S\pm P)}$; (5.2)

$[(1-\eta_H^2)\sim(\eta_H)]^{K(Z\pm S\pm P)}=1$ is called the logarithm of the unit circle.

Table is

(1) The various combination elements in the unit

have corresponding continuous and discontinuous, sparse and non-sparse, complete and incomplete spaces, positions, values, events in the range of unit.

(2) Each of the same-level group algebraic closed chains can accommodate the unity of each branch point, and the polynomial powers are expanded in the combination by a natural number positive integer infinite program.

(3) Based on the closed-chain combination of group algebra, the formed algebraic cluster power function and the circular logarithmic equation have integer change synchronism, avoiding the traditional mathematics "fixing a logarithm of a certain value (or constant)" cannot be eliminated" The residual number ε" realizes the integer zero error expansion, ensuring the smoothness and stability of the logarithmic equation and the limit value.

If the unitary circular logarithm theorem introduces the prime number theorem (PNT), for the condition of "sparse and non-sparse" of the prime distribution, the position and value of the distribution can still be determined within the unit logarithm of the unit, satisfying the Riemann conjecture. The requirement to determine the number of prime numbers before a certain value is known.

This unitary circular logarithm theorem, if introduced into the physical quantum theorem, ensures several aspects of superiority under the unitary conditions of quantum computing:

(1) Reciprocal and convertible interaction with positive and negative in each quantum system. Or solve the mystery of quantum computing.

(2) The entangled type calculation with uncertainty is converted into the calculation of relative certainty, and can explain the position, energy and the like of the long-distance transmission of each entangled particle in the wide area.

(3) It has the fields of algebra, geometry, numerical value, topology, probability, chaos, etc. that can be extended to any high dimensional dimension.

[Theorem 2], Tthe second norm-invariant theorem (reciprocal logarithm):

The reciprocal circular logarithm: "The average of the elements of the element is divided by the average of the total items of the elements, and the logarithm of the reciprocal circle is obtained."

heve：

$$(1-\eta^2)^{K(Z\pm S\pm P)}$$
$$=\{x_h\}^{K(Z\pm S\pm P)}/\{x_{0H}\}^{K(Z\pm S\pm P)}$$
$$=x_0^{K(Z\pm S-0)/t}D_0^{K(Z\pm S+0)/t}+x_0^{K(Z\pm S-1)/t}D_0^{K(Z\pm S+1)/t}+\cdots$$
$$+x_0^{K(Z\pm S-p)/t}D_0^{K(Z\pm S+p)/t}+\cdots+x_0^{K(Z\pm S-q)/t}D_0^{K(Z\pm S+q)/t}$$
$$\pm\{{}^{KS}\surd D\}^{+(Z\pm S+p)/t}\ ; \qquad (6.1)$$
$$(1-\eta^2)^{K(Z\pm S)}$$
$$=(1-\eta^2)^{K(Z\pm S-p)/t}+(1-\eta^2)^{K(Z\pm S+p)/t}+(1-\eta^2)^{K(Z\pm S\pm p)/t}$$
$$=\{1\}^{(Z/t)};\ \text{(odd function)} \qquad (6.2)$$

or: $$(1-\eta^2)^{K(Z\pm S)/t}=(1-\eta^2)^{K(Z\pm S+p)/t}+(1-\eta^2)^{K(Z\pm S-p)/t}$$
$$=\{1\}^{(Z/t)};\ \text{(even function)} \qquad (6.3)$$

and： $$(\eta)^{K(Z\pm S\pm P)}=\sum[(\eta_1)+\cdots+(\eta_p)]^{K(Z\pm S+P)}$$
$$+\sum[(\eta_2)\cdots+(\eta_q)]^{K(Z\pm S-P)}$$
$$=(\eta)^{K(Z\pm S+P)}+(\eta)^{K(Z\pm S-P)}$$
$$=\{1\}^{K(Z\pm S\pm P)}; \qquad (6.4)$$

Based on the multi-element multiplication of Lemma, the unitary logarithm of the unit is an arithmetic addition with a positive factor and an inverse factor, and the positive factor set and the inverse factor set form a set of factors, and vice versa. The symmetry of the mutuality of the unit logarithm of the unit.

Very easy to get: reversed by the logarithm of the circle,

heve：

$$(1-\eta^2)^{(Z/t)}=\prod(1-\eta^2)^{(Z/t)}=\sum(1-\eta^2)^{(Z/t)}\ ; \qquad (6.5)$$

The formula (7.1)~(7.5) is the process of obtaining the equilibrium set of the logarithmic factors of the positive and negative two types of circles after dividing the total elements of the group by the average value, and becomes the final equilibrium equation.

In particular, the reciprocity of the property of the logarithmic power function K=(+1,0,-1) breaks through the forbidden zone where the denominator of the Lobita rule is not zero. For example, the Riemann function is the sum of the reciprocals of prime numbers. When K = +1; S = +1, the harmonic progression is divergent. If the Riemann ζ function is the sum of the reciprocal of the prime numbers and then recounts, without losing its generality, the Riemann function K=-1; S=-1 is convergent, ensuring the stability of the convergence of the Riemann function. Expand.

[Theorem 3], The third norm-invariant theorem (homogeneous logarithm)

Proof: Algebraic closed chain algebraic clusters isomorphic reciprocal inversion

Let: $(1-\eta^2)^{K(Z/t)}=\sum[\{X_0\}/\{D_0\}]^{(Z/t)}$

have reciprocal inversion.

heve：

$$Ax^{K(Z\pm S\pm N-0)/t}+Bx^{K(Z\pm S\pm N-1)/t}+\cdots$$
$$+Px^{K(Z\pm S\pm N-p)/t}+\cdots+Qx^{K(Z\pm S\pm N-q)/t}+D^{K(Z\pm S\pm N+0)/t}$$
$$=[\{C_{(s+0)}x^{K(Z\pm S\pm N-0)/t}\bullet D_0^{K(Z\pm S\pm N+0)/t}\}$$
$$+\{C_{(s+1)}x^{K(Z\pm S\pm N-1)/t}\bullet D_0^{K(Z\pm S\pm N+1)/t}\}+\cdots$$
$$+\{C_{(s+p)}x^{K(Z\pm S\pm N-p)/t}\bullet D_0^{K(Z\pm S\pm N+p)/t}\}+\cdots$$
$$+\{C_{(s+q)}x^{K(Z\pm S\pm N-q)/t}\bullet D_0^{K(Z\pm S\pm N+q)/t}\}$$
$$\pm\{C_{(s+0)}{}^{KS}\surd D\}/\{D_0\}]^{K(Z\pm S\pm N+0)/t}$$
$$=[(1-\eta^2)^{K(Z\pm S\pm N+0)/t}+(1-\eta^2)^{K(Z\pm S\pm N+1)/t}+\cdots$$
$$+(1-\eta^2)^{K(Z\pm S\pm N+p)/t}+\cdots+(1-\eta^2)^{K(Z\pm S\pm N+q)/t}]$$
$$/\{X_0\pm D_0\}]^{(Z/t)}$$

$=(1-\eta^2)^{K(Z\pm S\pm N)/t}\{X_0\pm D_0\}]^{K(Z\pm S\pm N)/t}$; (7.1)

Obtained: circular logarithmic isomorphism: the isomorphism of regularized polynomials (including calculus equations)

$(1-\eta^2)^{K(Z\pm S\pm N)/t}=(1-\eta^2)^{K(Z\pm S\pm N+0)/t}$
$=\cdots=(1-\eta^2)^{K(Z\pm S\pm N+1)/t}=\cdots=(1-\eta^2)^{K(Z\pm S\pm N+p)/t}$; (7.2)

among them:

$(1-\eta^2)^{K(Z/t)}$
$=\{C_{(s+p)}x^{K(Z\pm S\pm N-p)/t}\cdot D_0^{K(Z\pm S\pm N+p)/t}\}$
$=\{x_0^{K(Z\pm S\pm N-p)/t}\cdot D_0^{K(Z\pm S\pm N+p)/t}\}$
$=\{x_0/D_0\}^{K(Z\pm S\pm N+p)/t}$; (7.3)

The isomorphic circular logarithm reflects the temporal algorithmic consistency of the polynomial inequality conversion equations of various (±P) combinations of algebraic clusters under the condition of regular regularization of equality and inequality. This makes any nonlinear problem convertible to a linear problem. The polynomial inequality conversion equation time algorithm has isomorphic uniformity.

where:

Z/t=+1, forward homeomorphic topology convergence process function, and finally a dot;

Z/t= 0, center point balance function;

Z/t=-1, the inverse boundary homeomorphic topology expansion function, and finally a circle;

[Theorem 4], the circular logarithm (relativistic construction) limit theorem

The polynomial or geometric space $(1-\eta^2)^{K(Z/t)}$ satisfies the multiplication of each level and converts it into a positive number and a reciprocal, which can be attributed to the isomorphic unity of a circular log-random stochastic topology and the unit stability of algebraic closed-chain Sexual limit, which is:

$(1-\eta^2)^{(Z/t)}=\Pi(1-\eta^2)^{(Z/t)}=\Sigma(1-\eta^2)^{(Z/t)}$; (8.1)

heve:

$(1-\eta^2)^{+(Z/t)}+(1-\eta^2)^{-(Z/t)}=1$; (8.2)

$(1-\eta^2)^{+(Z/t)}\cdot(1-\eta^2)^{-(Z/t)}=1$; (8.3)

Solution (8.2), (8.3) simultaneous equations, obtained: the logarithmic limit value of the stability, the critical value, and the boundary point.

$|(1-\eta^2)\sim(\eta)|^{K(Z/t)}$
$=(0,1/2,1)^{K(Z/t)}$
$=\{0,1/2,1,2\}^{K(Z/t)}$; (8.4)

when:

$|(1-\eta^2_{(r,\varphi,\theta,x,y,z)})\sim(\eta_{(r,,\varphi,\theta,x,y,z)})|^{K(Z/t)}$

heve:

$\eta_{(x,y,z)}=[0,1/2,1,2]^{K(Z/t)}$
(orthogonal coordinate system) (8.5)

or:

$\eta(r,\varphi,\theta)=[\theta_0\pm(0,\ \pi/4,\pi,2\pi)]^{K(Z/t)}$
(spherical coordinate system) (8.6)

when:

$|(1-\eta^2)\sim(\eta)|^{K(Z/t)}=(0,1/2,1)^{K(Z/t)}=\{0,1/2,1,2\}^{K(Z/t)}$;

At the limit, when applied to the Riemann conjecture, it can ensure the abnormal zero stability of the sum of the prime numbers of any positive and negative forms of the Riemannian function, which is $\{1/2\}^{K(Z/t)}$ at the critical line, which satisfies Riemann conjecture proved that the two conditions were required.

[Theorem 5], parallel/serial circular logarithm theorem

The composite hierarchical dynamic equation often consists of multi-level parallel equations with different elemental parameters to form a composite hierarchical dynamic equation. It is an important calculation method to clarify multi-level parallel equations. The stochastic decomposition of unit states based on group algebra closed chain becomes a parallel / serial polynomial equation, and the parallel/serial logarithm theorem is obtained.

There are: parallel/serial polynomial dynamic equation power function:

$(Z/t)=K(Z\pm S\pm(N_A+N_B+\cdots NP+\cdots Nq))/t$;
$(N_H)=(N_A+N_B+\ldots N_P+\ldots N_q))$;

Serial equation:

$\{D_{H0}\}^{(Z/t)}=\{(^{KS}\sqrt{}(D_i)\}^{(Z/t}$
$^{)}=(1/C_{(S\pm H)})\{D_A\cdot D_B\cdot\cdots\cdot D_P\cdot\cdots\cdot D_q\}^{(Z/t)}$,

Parallel equation:

$\{D_{H0}\}^{(Z/t)}=\{(\Sigma(D_i)\}^{(Z/t)}$
$=\Sigma(1/C_{(S\pm H)})\{D_A+D_B+\cdots+D_P+\cdots+D_Q\}^{(Z/t)}$,

Get the parallel/serial dynamic equation:

$\{x\pm D\}^{(Z/t)}=Ax^{K(Z\pm S\pm N-0)/t}+Bx^{K(Z\pm S\pm N-1)/t}+\cdots$
$+Px^{K(Z\pm S\pm N-p)/t}+\cdots+Qx^{K(Z\pm S\pm N-q)/t}+D^{K(Z\pm S\pm N+0)/t}$
$=(1-\eta_A^2)^{K(Z\pm S\pm N\pm A)/t}\{x_{0A}\pm D_{0A}\}^{K(Z\pm S\pm N\pm A)/t}$
$+(1-\eta_B^2)^{K(Z\pm S\pm N\pm B)/t}\{x_{0B}\pm D_{0B}\}^{K(Z\pm S\pm N\pm B)/t}+\cdots$
$+(1-\eta_P^2)^{K(Z\pm S\pm N\pm P)/t}\{x_{0P}\pm D_{0p}\}^{K(Z\pm S\pm N\pm P)/t}+\cdots$
$+(1-\eta_Q^2)^{K(Z\pm S\pm N\pm Q)/t}\{x_{0Q}\pm D_{0Q}\}^{K(Z\pm S\pm N\pm Q)/t}$
$=(1-\eta_A^2)^{K(Z\pm S\pm N\pm A)/t}\{0,2\}^{K(Z\pm S\pm N\pm A)/t}\{D_{0A}\}^{K(Z\pm S\pm N\pm A)/t}$
$+(1-\eta_B^2)^{K(Z\pm S\pm N\pm B)/t}\{0,2\}^{K(Z\pm S\pm N\pm B)/t}\{D_{0B}\}^{K(Z\pm S\pm N\pm B)/t}+\cdots$
$+(1-\eta_P^2)^{K(Z\pm S\pm N\pm P)/t}\{0,2\}^{K(Z\pm S\pm N\pm P)/t}\{D_{0P}\}^{K(Z\pm S\pm N\pm P)/t}+\cdots$
$+(1-\eta_Q^2)^{K(Z\pm S\pm N\pm Q)/t}\{0,2\}^{K(Z\pm S\pm N\pm Q)/t}\{D_{0Q}\}^{K(Z\pm S\pm N\pm Q)/t}$
$=(1-\eta^2)^{(Z/t)}\{0,2\}^{(Z/t)}\{D_0\}^{(Z/t)}$
$=(1-\eta^2)^{(Z/t)}[\{X_0\}^{(Z/t)}\pm\{D_0\}^{(Z/t)}]$
$=(1-\eta^2)^{-(Z/t)}\{X_0\}^{(Z/t)}\pm(1-\eta^2)^{+(Z/t)}\{D_0\}^{(Z/t)}$; (9.1)

$(1-\eta^2)^{(Z/t)}=(1-\eta^2)^{K(Z\pm S\pm N\pm[A+B+P+Q])/t}$
$=(1-\eta_A^2)^{K(Z\pm S\pm N\pm A)/t}+(1-\eta_B^2)^{K(Z\pm S\pm N\pm B)/t}+\cdots$
$+(1-\eta_P^2)^{K(Z\pm S\pm N\pm P)/t}+\cdots+(1-\eta_Q^2)^{K(Z\pm S\pm N\pm Q)/t}$; (9.2)

Algebraic closed-chain totals, sub-items, and branch terms have serial isoparagraphy with isomorphic compatibility of equations and inequalities:

$(1-\eta^2)^{(Z/t)}$

$= \Sigma\ {}^{KS}\sqrt{\{D_{A0} \bullet D_{B0} \bullet\cdots \bullet D_{P0} \bullet\cdots \bullet D_{Q0}\}^{(Z/t)}} / \{{}^{KS}\sqrt{(\prod D_{H0}\}^{(Z/t)}}$

$= \Sigma \{D_{A0}+D_{B0}+\cdots+ D_{P0}+\cdots+ D_{Q0}\}^{(Z/t)}/\{D_{H0}\}^{(Z/t)}$

$= (1-\eta_A^2)^{(Z/t)}+ (1-\eta_B^2)^{(Z/t)}+\cdots+(1-\eta_P^2)^{(Z/t)}+\cdots + (1-\eta_q^2)^{(Z/t)}$ (9.3)

or:

$(\eta)^{(Z/t)}$

$= (\eta_A)^{(Z/t)}+ (\eta_B)^{(Z/t)}+\cdots+(\eta_P)^{(Z/t)}+\cdots+ (\eta_q)^{(Z/t)}$; (9.4)

Each child has an interactive inversion:

$(1-\eta^2)^{K(Z\pm S)/t}=(1-\eta^2)^{K(Z-S)/t}\pm(1-\eta^{2\ K(Z+S)/t}$; (9.5)

Each sub-item has its own three-dimensional space

ates (the main

rotation equation):

$(1-\eta^2)^{K(Z\pm S)/t}$

$=(1-\eta_{[x]}^2)^{K(Z-S)/t}\mathbf{i}+(1-\eta_{[y]}^2)^{K(Z+S)/t}\mathbf{j}+(1-\eta_{[z]}^2)^{K(Z+S)/t}\mathbf{k}$ (9.6)

Each sub-item has its own three-dimensional Spherical Coordinates (spin equation):

$(1-\eta^2)^{K(Z\pm S)/t}$

$=(1-\eta_{[ZY]}^2)^{K(Z-S)/t}\mathbf{i}+(1-\eta_{[xz]}^2)^{K(Z+S)/t}\mathbf{j}+(1-\eta_{[xy]}^2)^{K(Z+S)/t}\mathbf{k}$; (9.7)

Among them: $(1-\eta^2)^{K(Z\pm S\pm[A+B+P+Q]\pm N)/t}$ has infinite expansion of polynomials and calculus of their respect.

The parallel serial logarithm theorem of Theorem 5 reflects that they are all arithmetically superposed by circular logarithmic factors, which is the main way to improve the performance of computer systems. Almost all high-performance computer systems, from SMP workstations and servers, CC-NUMA large servers, to supercomputer systems, are more or less parallel processing technologies. However, the introduction of traditional parallel processing technology also brings the defects of poor performance and poor programmability. Here, the parallel serial circular logarithm theorem integrates discrete parallel and entangled serial computing, making the time complexity directly equivalent to the computation time of the traditional processor.

The parallel algorithm-inclusive parallel/serial integrated system structure and software optimization technology are closely combined to create excellent conditions for the development of supercomputer theory.

6. Circular logarithm algorithm falsification Fermat-Wills inequality theorem

From the perspective of the development of mathematical history, the essence of the Fermat-Wills inequality theorem should be to solve the problem of how the inequality is transformed into a unified equilibrium equation. How to achieve parallel/serial inequality unification? The following is a reasonable proof of no proof by the exploration of Fermat-Wills inequality.

Proof: The power function is invariant, and the equations and inequalities are arbitrary integers or prime numbers that maintain the completeness and integrity equations.

Completeness: refers to the combination of two complete "if and only" groups as a complete integer group.

$\{A\}^{K(Z\pm S\pm P)}+\{B\}^{K(Z\pm S\pm P)}$

$=(1/2)\{A\pm B\}^{K(Z\pm S)}=\{C\}^{K(Z\pm S)}$

Integrity: refers to the combination of two integrity inequalities into a complete integer group.

$\{A\}^{K(Z\pm S\pm P)}+\{B\}^{K(Z\pm S\pm P)}$

$=(1-\eta^2)^{K(Z\pm S\pm P)}\ \{C\}^{K(Z\pm S\pm P)}$;

Where: {A}, {B}, {C} are all integers or prime expansions. (the same below) group set (completeness) inequality conversion equation sufficient proof.

Let: $\{A\}^{K(Z\pm S\pm P)}$, $\{B\}^{K(Z\pm S\pm P)}$, $\{C\}^{K(Z\pm S\pm P)}$ be an arbitrary complex of group algebra closed chain infinite (Z) elements Dimensions (±S), algebraic clusters (±P) combination levels and sets.

$\{X\}=\{A\}^{K(Z\pm S\pm P)}+\{B\}^{K(Z\pm S\pm P)}=(1/2)\{A\pm B\}^{K(Z\pm S)}$

Parallel combination of groups;

$\{X_0\}^{K(Z\pm S)}=\{{}^{KS}\sqrt{D_A}+{}^{KS}\sqrt{D_B}\}^{K(Z\pm S)}$; $\{D_0\}^{K(Z\pm S)}=(1/C_{(S+1)})\ (\Sigma X_i)^{K(Z\pm S\pm 1)}$;

$D_0 = D_{0A}+D_{0B}$;

Corresponding parameter combination

$\{C\}^{K(Z\pm S\pm P)}=(1/2)\{X\pm D\}^{K(Z\pm S)}$ means $\{A\}^{K(Z\pm S\pm P)}$ and $\{B\}^{K(Z\pm S\pm P)}$

Two groups have their own center points

$\{C_0\}^{K(Z\pm S\pm P)}=\{C_{A0}\}^{K(Z\pm S\pm P)}+\{C_{0B}\}^{K(Z\pm S\pm P)}$

a set of two center point averages;

$(1-\eta^2)^{K(Z\pm S)}=[\{X\}/\{D\}]^{K(Z\pm S)}=[\{X_0\}/\{D_0\}]^{K(Z\pm S)}$ central numerical function

topology change。

$(1-\eta_A^2)^{K(Z\pm S)}=[\{X\}/\{D_A\}]^{K(Z\pm S)}=[\{X_0\}/\{D_{0A}\}]^{K(Z\pm S)}$

central average numerical function topology change。

$(1-\eta_B^2)^{K(Z\pm S)}=[\{X_B\}/\{D_B\}]^{K(Z\pm S)}=[\{X_{0B}\}/\{D_{0B}\}]^{K(Z\pm S)}$ The central average numerical function is topologically changed.

After extracting the logarithm of the circle, making each level

$\{X_0\}^{K(Z\pm S)}=\{D_0\}^{K(Z\pm S)}=(1/C_{(S+1)})\ (\Sigma X_i)^{K(Z\pm S)}$;

regularization coefficient: $C_{(S+P)}=C_{(S-P)}$;

[6.1], The first type of proof retains the intermediate topology process:

Introduce the logarithm of the circle, select the relative comparison between their completeness and the final result, and the inequality becomes equation.

assume: $\{A\}^{K(Z\pm S\pm P)}+\{B\}^{K(Z\pm S\pm P)}=\{X\}^{K(Z\pm S\pm P)}$

$\{D_A\}^{K(Z\pm S\pm P)}+\{D_B\}^{K(Z\pm S\pm P)}=\{D\}^{K(Z\pm S\pm P)}$;

$\{D_A\}^{K(Z\pm S\pm P)}=\{A\}^{K(Z\pm S\pm P)}/[\{A\}^{K(Z\pm S\pm P)}+\{B\}^{K(Z\pm S\pm P)}]$;

$\{D_B\}^{K(Z\pm S\pm P)}=\{B\}^{K(Z\pm S\pm P)}/[\{A\}^{K(Z\pm S\pm P)}+\{B\}^{K(Z\pm S\pm P)}]$;

$(1-\eta^2)^{K(Z\pm S\pm P)}=\{X\}^{K(Z\pm S\pm P)}/\{D\}^{K(Z\pm S\pm P)}$;

$(1-\eta_A^2)^{K(Z\pm S\pm P)}=\{D_A\}^{K(Z\pm S\pm P)}/\{D\}^{K(Z\pm S\pm P)}$;

$(1-\eta_B^2)^{K(Z\pm S\pm P)}=\{D_B\}^{K(Z\pm S\pm P)}/\{D\}^{K(Z\pm S\pm P)}$;

heve:

$$
\begin{aligned}
&\{X\pm D\}^{K(Z\pm S)}=\{AX^{K(Z\pm S\pm 0)}+BX^{K(Z\pm S\pm 1)}+\cdots+PX^{K(Z\pm S\pm P)}+\cdots+QX^{K(Z\pm S\pm q)}\pm D_A\}\\
&+\{AX^{K(Z\pm S\pm 0)}+BX^{K(Z\pm S\pm 1)}+\cdots+PX^{K(Z\pm S\pm P}+\cdots+QX^{K(Z\pm S\pm q)}\pm D_B]\\
&=\{X_A^{K(Z\pm S\pm 0)}+C_{(S-1)}X_A^{K(Z\pm S-1)}D_{0A}^{K(Z\pm S\pm q+1)}\\
&+C_{(S-p)}X_A^{K(Z\pm S-P)}D_{0A}^{K(Z\pm S+p)}\\
&+C_{(S-q)}X_A^{K(Z\pm S-q)}D_{0A}^{K(Z\pm S+q)}\pm D_A\}\\
&+\{X_B^{K(Z\pm S\pm 0)}+C_{(S-1)}X_B^{K(Z\pm S-1)}D_{0B}^{K(Z\pm S\pm q+1)}\\
&+C_{(S-p)}X_B^{K(Z\pm S-P)}D_{0B}^{K(Z\pm S+p)}\\
&+C_{(S-q)}X_B^{K(Z\pm S-q)}D_{0B}^{K(Z\pm S+q)}\pm D_B\}\\
&=\{X_{0A}^{K(Z\pm S\pm 0)}+X_{0A}^{K(Z\pm S-1)}D_{0A}^{K(Z\pm S\pm q+1)}\\
&+X_{0A}^{K(Z\pm S-P)}D_{0A}^{K(Z\pm S+p)}\\
&+X_{0A}^{K(Z\pm S-q)}D_{0A}^{K(Z\pm S+q)}\pm D_{0A}^{K(Z\pm S+0)}\}\\
&+\{X_{0B}^{K(Z\pm S\pm 0)}+X_{0B}^{K(Z\pm S-1)}D_{0B}^{K(Z\pm S\pm q+1)}\\
&+X_{0B}^{K(Z\pm S-P)}D_{0B}^{K(Z\pm S+p)}\\
&+X_{0B}^{K(Z\pm S-q)}D_{0B}^{K(Z\pm S+q)}\pm D_{0B}^{K(Z\pm S+0)}\}\\
&=\{(1-\eta_A^2)^{K(Z\pm S)}\{X_0\pm D_{0A}\}^{K(Z\pm S)}\}\\
&+\{(1-\eta_B^2)^{K(Z\pm S)}\{X_0\pm D_{0B}\}^{K(Z\pm S)}\}\\
&=(1-\eta^2)^{K(Z\pm S)}\{X_0\pm D_0\}^{K(Z\pm S)}\\
&=(1-\eta^2)^{K(Z\pm S)}[\{X_0\}^{K(Z\pm S)}\pm\{D_0\}^{K(Z\pm S)}]\\
&=(1-\eta^2)^{K(Z\pm S)}\{2\}^{K(Z\pm S)}\cdot\{D_0\}^{K(Z\pm S)};
\end{aligned}
\qquad (10.1)
$$

$\{A\}^{K(Z\pm S\pm P)}+\{B\}^{K(Z\pm S\pm P)}$

$=(1-\eta^2)^{K(Z\pm S)}[2]\cdot\{C\}^{K(Z\pm S\pm P)}$; (10.2)

$(1-\eta^2)^{K(Z\pm S\pm P)}$

$=(1-\eta_A^2)^{K(Z\pm S\pm P)}+(1-\eta_B^2)^{K(Z\pm S\pm P)}$; (10.3)

among them:

$\{2\}/\{2\}^{K(Z\pm S\pm P)}$

$=[\{A\}^{K(Z\pm S\pm P)}+\{B\}^{K(Z\pm S\pm P)}]/\{A+B\}^{K(Z\pm S\pm P)}$

[6.2], The second type proves the direct topology process

Introduce the logarithm of the circle and directly select the relative comparison of their final results. The inequality becomes the equation.

$(1-\eta^2)^{(Z/t)}\sim(\eta)^{(Z/t)}=[(A+B)-C]/[(A+B)+C]$

$=[(A^{K(Z\pm S\pm P)}+B^{K(Z\pm S\pm P)})-C^{K(Z\pm S\pm P)}]$

$/[(A^{K(Z\pm S\pm P)}+B^{K(Z\pm S\pm P)})+C^{K(Z\pm S\pm P)}]$; (10.4)

heve:

$\{A\}^{K(Z\pm S\pm P)}+\{B\}^{K(Z\pm S\pm P)}=(1-\eta^2)^{K(Z\pm S\pm P)}\{C\}^{K(Z\pm S\pm P)}$; (10.5)

among them:

$\{A\}^{K(Z\pm S\pm P)}=(1-\eta_A^2)^{K(Z\pm S\pm P)}\{C\}^{K(Z\pm S\pm P)}$; (10.6)

$\{B\}^{K(Z\pm S\pm P)}=(1-\eta_B^2)^{K(Z\pm S\pm P)}\{C\}^{K(Z\pm S\pm P)}$; (10.7)

$(1-\eta^2)^{K(Z\pm S\pm P)}=(1-\eta_A^2)^{K(Z\pm S\pm P)}+(1-\eta_B^2)^{K(Z\pm S\pm P)}$; (10.8)

[6.3], Proof of the type of reverse inequality

If $(1-\eta^2)^{K(Z\pm S\pm P)}$ and $\{C\}^{K(Z\pm S\pm P)}$ are known, look for $\{A\}^{K(Z\pm S\pm P)}$ and $\{B\}^{K(Z\pm S\pm P)}$ becomes a problem of inverse inequality. Under the condition of the reciprocity theorem.

$(1-\eta_A^2)^{K(Z\pm S\pm P)}+(1-\eta_B^2)^{K(Z\pm S\pm P)}=\{1\}^{K(Z\pm S\pm P)}$; (11.1)

or:

$(1-\eta_A^2)^{K(Z\pm S+P)}\cdot(1-\eta_B^2)^{K(Z\pm S-P)}$

$=\{1\}^{K(Z\pm S\pm P)}$; (11.2)

heve:

$(1-\eta_A^2)^{K(Z\pm S+P)}\rightarrow 0$; $(1-\eta_B^2)^{K(Z\pm S-P)}\rightarrow 1$;

the opposite is also true;

How to determine the value of {A} and {B} deterministically?

(6.3.1) Knowing the {A} and {B} composition rules and {C}K (Z±S±P),

Assume:

$\{C\}^{K(Z\pm S\pm P)}=(1-\eta_{AB}^2)^{K(Z\pm S\pm P)}[\{A\}+\{B\}]^{K(Z\pm S\pm P)}$

Presence $(\eta_{AB})^{K(Z\pm S+P)}=(\eta_{BA})^{K(Z\pm S-P)}$

Where: the combination of integrity is {2}・{C}, and the completeness is {1}・{C} heve: $(\eta_{AB})^{K(Z\pm S\pm P)}=(\eta_{BA})^{K(Z\pm S\pm P)}=[C_{0AB}-\{A\}]/C_{0AB}=[\{B\}-C_{0AB}]/C_{0AB}$

get:

$$\{A\}^{K(Z\pm S\pm P)}=(1-\eta_A{}^2)^{K(Z\pm S+P)}C_{0AB} = (1-\eta_{AB})^{K(Z\pm S+P)}C_{0AB}; \quad (11.3)$$

$$\{B\}^{K(Z\pm S\pm P)}=(1-\eta_B{}^2)^{K(Z\pm S-P)}C_{0AB} = (1+\eta_{BA})^{K(Z\pm S-P)}C_{0AB}; \quad (11.4)$$

(6.3.2), I do not know $\{A\}^{K(Z\pm S\pm P)}$ and $\{B\}^{K(Z\pm S\pm P)}$ composition rules and $\{C\}^{K(Z\pm S\pm P)}$, $\eta_{AB})^{K(Z\pm S+P)}=(\eta_{BA})^{K(Z\pm S-P)}$

From an information point of view, it is a security measure that is temporarily confidential. Cryptography involving security privacy.

The logarithm of the circle reflects the common rules of information. For different unknowns, only rely on scientific experiments or multiple guesses. Scientific experiments are indispensable means for human beings to explore nature. For this reason, it is necessary to improve the design and improvement of computer functions in the future.

7. Physics, Mathematical Verification and Engineering Application of Fermat's Theorem

7.1. The reliability of the theory of circular logarithm is proved by the physical experiment "angel particle"

In July 2017, the Zhang Shouyi team discovered the edge current with half quantum conductance in the superconducting-quantum anomalous Hall platform, which is in good agreement with the theoretically predicted chiral Mayorana fermion. This is the first one.

The evidence of the Mayalana measurement. It is another milestone discovery after the “God particle”, “neutrino” and “gravity”. "Angel particles" are the only positive and negative particles in the fermion. It coincides with the findings of other particles, which proves the reliability and authenticity of the theory of circular logarithm.

heve:

$$E=(1-\eta^2)^Z MC^2 \quad (12.1)$$

$$\{0\}^Z \leqslant (1-\eta^2)^Z \leqslant \{1\}^Z; \quad (12.2)$$

Where: $Z=K(Z\pm S\pm N)/t$; $(MC2)$ unchanged,

Under equilibrium conditions: $\{2\}^Z$ is the quantum bit entanglement information.

Such as angel particles:

$(0)\leqslant(1-\eta^2)^Z\leqslant(1)$: $(K=+1,0,-1)$:

Angel particle positive particle:

$(0)\leqslant(1-\eta^2)^{+Z}\leqslant(1/2)$:

Angel particle antiparticle:

$(1/2)\leqslant(1-\eta^2)^{-Z}\leqslant(1)$:

The inequalities and equations are unified by the logarithm of the circle, so that any integers and primes remain complete and complete. Also obtained a lot of physical experiments.

7.2. In theory, there are series of results verification:

The essence of Fermat's theorem inequality problem is the uncertainty problem. Wiles proves its uncertainty, so the Fermat's theorem is established. Corresponding to the microscopic uncertainty principle of Heisenberg's quantum mechanics, it is said that "position and kinetic energy cannot be determined simultaneously" in microscopic quantum. That is to say, once the quantum position is determined, the kinetic energy cannot be determined, and vice versa. However, under the multi-particle, multi-element and multi-space, Yang Zhenning-Mills wrote a "normative field", trying to achieve the great unity of natural forces. This uncertainty is even more difficult. The mathematicians have requested that "there is no calculation of the quality element content." "Resolved, called the seven mathematical problems of the 21st century.

The author believes that many of today's mathematical problems are essentially a problem. These problems are all related to Fermat's theorem inequality problem. The circular log-polynomial method can be used to meet the requirements.

In 2018, Wang Yiping's theory of circular logarithm was published in the American Journal of Mathematics and Statistics. There are "Circular Logarithm and Riemann Function" (JMSS 2018/1); "Circular Logarithm and Gauge Field" (JMSS 2018/2); "Circular Logarithm and NS Equations and Applications" (JMSS 2018/5); More than 6 articles such as the complete logarithm and P-NP complete problem (JMSS 2018/9).

7.3. There are a series of application verifications on the project.

(7.3.1) Wang Yiping, Quzhou City, Zhejiang Province, “Two-way Vortex Vacuum Energy Engine”, proposed the four-stroke working system of the engine to create asymmetric energy and reform the working principle of the traditional engine. In 2014-2015, he obtained two Chinese national invention patents.

(7.3.2) Sun Chunwu (Wang Yiping Research Team) of Yangzhou City, Jiangsu Province, “Eccentric Rotating Engine”, proposed the application of Earth's gravity energy, determined the relationship between spin position and kinetic energy, and converted uncertainty into certainty. In 2017, I applied for a national invention patent. The prototype has been successfully manufactured.

(7.3.3) The “No Block Hydropower Station” in cooperation between Wang Yiping of Zhangzhou City, Zhejiang Province and Yangjingshan of Anyang, Henan Province, uses a variety of gravity energy to generate water. It has been mass produced.

(7.3.4) Xu Wenyu, “Small Sub-Magnetic Energy Superimposed Vector Power Generation Device” of Shenzhen City, Guangdong Province, won the national invention patent. The prototype has been successfully manufactured.

These inventions all apply the inequality to the equation principle. It has a traditional engine with environmental protection and energy conservation, and a power mechanical device that uses gravity field and magnetic field as energy.

Therefore, Fermat's theorem is transformed into an equation, which proves that any inequality can be expanded by a logarithm to an integer of the integer solution. Has significant mathematical and engineering application value.

8. Looking into the digital world of the 21st century

Quantum computing is the translation of our macro information into the mechanical quantities of particles, such as the first-order, second-order equations of polynomials and even any higher-order angular momentum, momentum, and force. Then through the particle motion, the mechanical quantity is converted into macro information. Due to their homogeneity and reflexivity, they appear in pairs under certain conditions and can disappear in pairs. The former is applied to information transfer, which is to maintain the parameters of the operation, the latter can apply the conversion, that is, to do arithmetic four operations.

The concept of circular log-blockchain-quantum computing extension attempts to be able to deal with existing, whole-world operations. But one of the reasons is not easy to implement. One of the reasons is that people want to "entangle" all the qubits, let them maintain the quantum state for a long time. In mathematics, the integer expansion, that is, the first need to deal with the Hodge conjecture, NP- P complete problems, Riemann's conjecture and other problems, as well as a series of problems involving the computer itself, production and so on. It can be seen that creating a quantum computer is a comprehensive problem.

8.1, The round logarithm - the combination of blockchain

From the point of view of the unitary topological blockchain: the logarithm of the circle can reflect the superposition state and the state of the parallel/serial superposition of the particles, and the arbitrary finite Z=K (Z±S±P) topological quantum changes of the respective infinite group spaces.

when: S = +p, boundary (maximum value) respectively, multidimensional algebra - geometric space converges to the center point of each and the homeomorphic inner circle (infinitesimal limit);

S = -p, conversely, spread by the (minimum value) boundary to the outer circular boundary condition (the limit of infinity);

S = ± p, is the multi-dimensional algebra - geometric space and the common inner geometric center point or boundary point, line, surface, body, multi-group aggregate limit.

Also known as singularity, critical point, sudden point, conversion.

heve：

$$W=(1-\eta^2)^{K(Z\pm S\pm N)/t}\,W_0\,; \qquad (13.1)$$

$$\{0\}^{K(Z\pm S\pm N)/t}\leqslant(1-\eta^2)^{K(Z\pm S\pm N)/t}\leqslant\{1\}^{K(Z\pm S\pm N)/t}; \qquad (13.2)$$

$$(1-\eta^2)^{K(Z\pm S\pm N)/t}=(1-\eta^2)^{K(Z+S+N)/t}+(1-\eta^2)^{K(Z-S-N)/t}; \qquad (13.3)$$

$$(1-\eta^2)^{K(Z\pm S\pm N)/t}=(1-\eta^2)^{K(Z\pm 0)/t}+(1-\eta^2)^{K(Z\pm 1)/t}\cdots+(1-\eta^2)^{K(Z\pm q)/t}; \qquad (13.4)$$

for serial:

$$(\eta)^{K(Z)/t}=(\eta)^{K(Z\pm 0)/t}+(\eta)^{K(Z\pm 1)/t}+\cdots+(\eta)^{K(Z\pm p)/t}+\cdots+(\eta)^{K(Z\pm q)/t}=(0\sim 1); \qquad (13.5)$$

for parallel:

$$(\eta)^{K(Z)/t}=(\eta_1)^{K(Z\pm p)/t}+(\eta_2)^{K(Z\pm p)/t}+\cdots+(\eta_p)^{K(Z\pm p)/t}+\cdots+(\eta_q)^{(Z\pm p)/t}=1; \qquad (13.6)$$

Combine and write a logarithmic equation: Combine and write a logarithmic quation:

$$(\eta)^{K(Z)/t}=(\eta_1)^{K(Z\pm 0)/t}+(\eta_2)^{K(Z\pm 1)/t}+\cdots+(\eta_p)^{K(Z\pm p)/t}+\cdots+(\eta_q)^{K(Z\pm q)/t}=(0\sim 1); \qquad (13.7)$$

The combination of equations (13.1)~(13.7) becomes a good time algorithm for random, topological quantum computing.

Among them: W, W_0 means any event, phenomenon, mechanics, space, value, etc. before and after. $(1-\eta^2)^{K(Z\pm S\pm N)/t}$: Topological and dynamical rules of any finite hierarchy, ensuring that any inequality is converted to an equation.

Thus, the power function of the circular log-block chain is in integer form, which satisfies the unity, reciprocity, isomorphism and stability of quantum computing. Where the logarithm of the circle is expanded, such as

(1), To meet the Decentralization: is a common rule acceptable to everyone.

(2) It satisfies the Security, has the same form of superposition state branch state, and the structure of different combinations protects the performance of the topological quantum.

(3), Satisfying the Scability. The data collection, processing data, and data modeling in data mining can be attributed to the infinite expansion between [0~1] without specific objects.

8.2、Bug measurement and control

However, although the existing blockchain is powerful, there are often bugs in branch computing, which hurts the reliability of the blockchain system. It is necessary to thoroughly understand the blockchain characteristics to design effective tools to prevent, control, and Relieve bugs. Collect valid data sets [11] through some of the following deployments.

(1) Retrieving representative blockchain items;

(2) Retrieving Bug reports and related information;

(3) Identifying the main language of the project;

(4) Calculating the duration of Bug repair;

After digitization, the blockchain and the bug have different polynomial features: the dimension (S), the coefficient {B 或 P}, the parameter $\{\sqrt[KS]{D}\}$, the average $\{D_0\}$, the logarithm of the circle $(1-\eta^2) \sim (\eta) = \{\sqrt[KS]{D}\}/\{D_0\}$. Make circle log-machine identification distinguishable and quickly corrected.

Through the calculation of the logarithmic integrity of the circle, the identification of zero error expansion is realized, which becomes the control and detection technology of the circular log-block chain.

8.3. Circular Logarithm-Complementary Compatibility and Expansion of Blockchain

The logarithm of a circle can represent the symbol of any thing by numbers, perform arithmetic four operations, and establish an algebra-geometric polynomial with "nothing." Since its inception, the blockchain has been described as a powerful technology with “all-powerfulness”.

In November 2018, “Digital” CEO Dai Weiguo pointed out “Digital Economy and Trusted World Research – Talking from Blockchain Infrastructure”, pointing out the shortcomings of the traditional digital economy, changing all of this requires a credible infrastructure. In accordance with the rules that are recognized by the public and can not be falsified - "blockchain infrastructure." It is necessary to write code, think about the underlying logic of the blockchain, and crack the impossible triangles. The impossible triangle is the blockchain network. No matter which consensus mechanism is adopted to determine the generation method of the new block, it is difficult to take into account both the scalability (Scability), the security (Security) and the decentralization. Three requirements. Although the American Massachusetts Institute of Technology Mikali proposed a program called Alogoran, it is a combination of "algorithm (algorithm)" and "random (random)". Limitedly used in the field of financial economy and network information.

Conclusion

Hilbert described Fermat's theorem as "a swan that will lay golden eggs" more than a hundred years ago. If you want to say why Fermat's theorem is so important in the history of mathematics, Wiles's sentence can be said: "Is it a good idea to judge whether a mathematical problem is good? The standard depends on whether it can produce new mathematics, not the problem itself.

Through the exploration of Fermat-Wills' theorem, we find the reciprocal mean, the unity of the difference and compatibility between the inequality and the equation, and establish the logarithm-circle logarithm with the dimensionless elliptic function as the base. The arithmetic implementation of the "zero error" expansion of integers and the dimensionless quantity without specific element content between [0 and 1] integrates inequalities with equations, or entangled and discrete types. A new, reliable, concise, and universal mathematical system was born.

The inequalities and equations are unified by the logarithm of the circle, so that any integers and primes remain complete and complete. Get a lot of physical experiments. On the quantum computer, the quantum bit entanglement problem can be solved smoothly. Convincingly falsified the Fermat-Wills inequality theorem.

It can be said that all the scientific problem bottlenecks in today's world, many problems are concentrated on the falsification of Fermat's theorem, only the round logarithm becomes the breakthrough of cracking. The perfect combination of circular log-blockchains can transform any random topology-probability event into a number and become a log-to-block chain. In addition to the performance of quantum computers in the future. There is no more secret in the world.

With the popularization and application of the circular log-blockchain, many mathematical formulas will be summarized as "four Latin letters". Surprisingly found: a simple formula is actually self-consistent to accommodate too much connotation of the mathematics building, reflecting the highest realm of "avenue to Jane". (Finish)

About the Author:

Wang Hongxuan, born in 2001, male journal of Jiangshan Experimental High School International Journal (JMSS) published three papers based on the

four-color theorem of circular logarithm-multivariate analysis (first author) "solving arbitrary high-dimensional polynomial equations", Received 8 patents including ultra-high pressure vane water pump of Chinese invention patent.

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Professor, 1961 Graduated from Zhejiang University Zhejiang Quzhou City Association of Old Scientists.

China Qianjiang Institute of Mathematics and Power Engineering Researcher Senior Engineer.

Engaged in basic mathematics and power mechanical engineering research.

The "reciprocal mean" was found, and a logarithmic equation of dimensionless quantities was established, which is widely used in cardinal and various scientific fields.

Internationally published papers include "Riemann function and relativistic structure", "P-NP complete problem and relativistic structure" and more than 10 articles, He obtained 6 items including the Chinese invention patent "Two-way scroll internal cooling engine" and "Two-way scroll hydrogen engine".

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